

Software Driven Readout Unit

B. Jost & N. Neufeld



Readout Unit (RU)

- Receives data from one or more front end links and assembles the fragments belonging to each event
- Complete fragments are then sent out to the readout Network



Requirements for RU

- Handle up to 4 input links at ~1 Gb/s each
- Write out data to the readout network, at a maximum rate of ~ 1 Gb/s (e.g.: Gigabit-Ethernet, Myrinet)
- Provide some buffering for short network congestions
- Eventually raise a throttle signal when unable to evacuate data



A proposal for a RU based on the IBM NP4GS3

- The IBM NP4GS3 is a network processor designed for wirespeed switching/routing and frame manipulation
- It can handle over 4.5 Million packets per second on 4 1-Gigabit full duplex links
- It is fully software programmable

IBM innovation *speeds* communication networks





Note:

NP4GS3-based System Architecture



CERN. EP



IBM NP4GS3 features

- 16 Special Purpose Processors @133 MHz each with 2 Hardware threads for Datahandling
- Always 2 processors share a set of coprocessors for all important tasks (data moving, table look-up, check-summing, etc.)
- High-speed interface to a switch (DASL)

- 128 kByte of storage on the input (ingress) side
- Up to 64 MByte on the output (egress) side
- Embedded PowerPC for remote management, monitoring, debugging and PCI access

LHCb NP4GS3 internal Architecture



EPC Block-Diagram



CERN. EP



RU Implementation using NPs



9



FEM Implementation using NPs

10









Software Environment

- I BM provides a rich software environment for the processor:
 - Assembler, Emulator and Debugger and Network traffic simulator
 - See Demo later...
- There exists also a hardware evaluation kit
 - The kit comprises NP, Gigabit Ethernet Ports, switch chip carrier, connector to Remote Debugger etc.
 - The Reference Evaluation Kit is a fully functional RU prototype (in principle)



Dataformat

- For technical reasons it is advantageous to have the blocks in a data fragment 64-byte aligned
- Data blocks consist of:
 - a header: event number, source id, offset of data block, offset of error block, and some others (64 Bytes total)
 - a data block (possibly empty)
 - a data padding to 64 bytes
 - an error block (possibly empty)
 - an error padding to 64 bytes







Data on the ingress side





(Sub)-Eventbuilding

- The original headers of the individual fragments are stripped of, a new header of the (sub-)event is created
- Data blocks are concatenated in the order of their source
- Error blocks are concatenated in the order of their source and appended to the data block
- The paddings have to removed from the data and error blocks
- The concatenated block have again to be padded to 64 Bytes, to preserve the original structure



Concatenation in Ingress



16







Performance of Event Building Code

- The first running version of the code takes 4.5 µs / fragment (for a single thread out of 30 general purpose threads!)
- Single thread execution time can easily be cut by a factor 30%
- All main important features of the processor understood



Advantages of Software RU

- Simple board with very few components: NP4GS3s, 5xPhy, (PC-)Memory-Chips
- Ready-to-buy evaluation board to implement any desired prototype
- "Full" flexibility
 - e.g. destination assignment as function of key word in header
 - traffic shaping
 - etc.



Next Steps

- Evaluate different software versions for Multiplexer, RU
- I mprove Running Speed of code
- Acquire a Reference Platform from IBM
- Study RU functionality (is a 'discrete' prototype)